

Scheduling of the inbound and outbound trucks at a Cross-Docking platform considering mutualization of docks

Natalia Cordoba^a , Alejandro Montoya^{a*} 

^aUniversidad EAFIT, Departamento de Ingeniería de Producción, Medellín, Colombia

*jmonto36@eafit.edu.co

Abstract

Paper aims: This paper studies a new Cross-Docking (CD) scheduling problem inspired by the trend of resource mutualization at CD operations where a company owning a CD platform can lease its docks to other companies.

Originality: This work pioneers the studies of the resource mutualization at CD operations solving at scheduling problem, where there is a set of requirements for leasing the loading and unloading docks for a period, and the company owning the docks can accept or decline each requirement. Considering that the availability of the docks is linked to scheduling the trucks at the company's loading and unloading docks, the decision to accept the requirements must be made along with the decisions for scheduling the trucks.

Research method: To solve this problem, this work proposes a Mixed-Integer Linear Programming (MILP) formulation.

Main findings: Using a commercial solver, the MILP formulation can optimally solve 65/128 instances based on a real operation. For the remaining instances, the MILP formulation obtains an average optimized gap of 8.95%.

Implications for theory and practice: The MILP formulation obtains acceptable results. Moreover, we found that all the instances accepted the requirements showing that it is economically interesting to lease docks.

Keywords

Cross docking platform. Mutualization. Scheduling problem. MILP formulation.

How to cite this article: Cordoba, N., & Montoya, A. (2022). Scheduling of the inbound and outbound trucks at a cross-docking platform considering mutualization of docks. *Production*, 32, e20210091. <https://doi.org/10.1590/0103-6513.20210091>

Received: July 21, 2021; Accepted: Nov. 24, 2021.

1. Introduction

Cross docking (CD) is a logistics distribution strategy used by many companies for arriving to clients faster; hence, being more competitive. It consists in transferring products from inbound trucks to outbound ones (van Belle et al., 2012) without storing or with a short one. Moreover, CD is used, for instance, at retailers for obtaining a last-mile distribution (Boysen & Fliedner, 2010). This model is used the companies such as Grupo Familia or Sodimac, which allows them to have a reduction in transportation costs (Siguenza, 2016).

Within traditional logistics, there are five standard processes in storing: reception, sorting, storing, recovery, and delivery (Boloori Arabani et al., 2010). When using a CD platform, it is expected to lower the storing and recovery times; therefore, eliminating their associated costs. To achieve these benefits at a CD platform, different activity scheduling problems need to be solved. One of the most common ones is determining the order at which the inbound and outbound trucks need to be served, as well as the assignation of the trucks at the docks. This decision is made aiming a synchronization between both inbound and outbound trucks that allows the reduction of waiting times for the trucks, their unnecessary displacement, or their temporary storing (Boysen et al., 2010).



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Given to current challenges in the companies, they are looking for alternatives to be more competitive. Within the alternative, there is the mutualization of resources for reducing operational and investment costs. Car manufactures are among some of the examples for mutualization. For instance, the alliance Renault-Nissan wherein the negotiation with the shipping companies is carried out conjunctively. This allows them to distribute their cars globally at better rates (Automotive Logistics, 2010). Another example comes from the principality of Monaco where policymakers created an urban platform for all of the distributors to the principality for complying with vehicle mobility restrictions within the city (Paché, 2008).

Following current trends on mutualization of logistics resources, this work solves a truck-scheduling optimization problem for a CD platform that mutualizes its docks with other allied companies. This problem is inspired by a real problem where a company owning a CD platform has a set of requirements for leasing its loading and unloading docks for a determined period to allied companies. Considering that the availability of the docks depends on the schedule of the inbound and outbound trucks belonging to the company, the decision of whether requirements are accepted or not need to be made along the truck scheduling decisions. The latter consists in defining the order in which the inbound and outbound trucks are served, as well as their assignation to a loading or unloading dock. The decisions for accepting requirements and scheduling the company's trucks need to be made in a way that minimize the total cost, which is defined as the truck-waiting costs minus the benefit that can be received with taking in requirements. In this problem, it is assumed that all the trucks have an arrival time and that all the inbound and outbound docks are mutually exclusive.

Regarding the objective function, it is important to clarify as follows. Even though within the objective function there is the income for accepting requirements, for the real operation that inspired this problem, the waiting cost is significantly higher than the income for accepting requirements. These incomes have to be interpreted more as a subsidy on the operational cost. Therefore, this work opted for minimizing the difference between the costs and the incomes. Furthermore, within the operation that inspired this problem, the costs associated with the platform are fixed (regardless of the scheduling for the trucks); consequently, the only cost that needs to be accounted for is the waiting cost associated with the trucks, which is representative.

For solving this problem, a Mixed-Integer Linear Programming (MILP) formulation is proposed. To evaluate the MILP formulation, we generated 128 instances based on a real operation from a retail company. Despite the instances coming from an industrial scale, our MILP formulation was capable of optimally solving 65/128 instances on a commercial solver with an average time of 5 minutes. For the remaining instances, our MILP formulation obtained an average optimally gap of 8.95% within a maximum time of 180 minutes. In practical terms, for the daily scheduling of a CD operation, this result is acceptable. The company that inspired this problem can run the MILP formulation every day, early in the morning, with a maximum limit time of 180 minutes; finding an optimal solution in half of the cases before the time limit or having a solution with an optimally gap relatively low within this time.

The rest of this paper is organized as follows: Section 2 presents a literature review associated to the problem at hand. Section 3 formally describes the problem and introduces the MILP formulation. Section 4 presents the computational study. Finally, Section 5 shows the conclusions.

2. Literature review

Optimization problems associated with CD platform operations have been studied by many authors in the literature. According to Buijs et al. (2014), these problems can be classified into strategical, tactical, and operational problems depending on the decisions that are made. Among the strategical decisions, there are those related to the design of the distribution network that integrates the operations of the CD platforms or the design of the CD platform's inner operation. Among the tactical decisions, there is the planning of the capacity in terms of the estimation of the needed resources (i.e., personnel and equipment) for handling an expected volume of goods. Finally, among the operational decisions, there are the programming of the resources within the platform, the assignment of the trucks to the docks, the order for tending to the trucks, among others (Buijs et al., 2014). Considering that the problem at hand is classified as an operational problem, the literary review will focus on the operational problems on truck scheduling. Moreover, considering that the novelty of this study is connected to the mutualization of the CD docks, some logistics-operation-mutualization cases will also be reviewed.

Hereafter, we will review the works that focus on operational problems, which are closest to this study. Yu & Egbelu (2008) studied the problem of finding the sequence of the trucks and their assignment to the docks in a way that minimize the makespan CD's operation. For solving this problem, the authors propose three different approaches, a MILP formulation, an exhaustive enumeration, and a heuristic method. The latter is used for solving medium and large-size problems. Conversely, Bozer & Carlo (2008) studied a problem where a decision

on which truck is served at which dock is made in a way that minimize the distances for moving the goods inside the platform, considering the rectilinear movements from the inbound trucks till the outbound trucks. For solving this problem, the authors propose a MILP formulation for the Quadratic Assignment Problem (QAP) with rectilinear distances and a Simulated Annealing heuristic. Boysen et al. (2010) presented a truck scheduling problem for minimizing the activity times considering the exact moment where the deliveries need to be made for the different products. For solving this problem, the authors use a decomposition method, which uses dynamic programming and some heuristics. Finally, Molavi et al. (2018) presented a truck scheduling problem where the delivery date of the products is the constraint. This problem seeks to minimize the total costs associated with fines and late deliveries. For solving this problem, the authors propose a hybrid genetic algorithm.

In recent years, different studies focusing on collaborative logistics have been developed. Furthermore, many analyze how the companies can work alongside their competitors (i.e., horizontal strategy), their providers, and their clients (i.e., vertical strategy) (Cleophas et al., 2019). From the literature review, different cases of collaborative logistics can be found. For instance, Blanquart & Carbone (2010) analyzed 7 collaborative logistics cases, four related to production routes for mass consumption, one related to product conception, and 2 related to reverse logistics. The authors classify the cases into 4 families considering if the collaborators belong or not to the same supply chain, and if their services or products are complementary. Another example of mutualization in logistics is the case analyzed by Morana (2010), wherein the French press mutualizes the urban distribution routes, so every customer is reached. Particularly, within the context of storing and transport, there are certain mutualization studies. Kreng & Chen (2008) created a collaborative system in which several companies must produce the good in a way that they simultaneously arrive in a distribution center for being able to perform a joint distribution; hence, lowering costs. The same year, Paché (2008) presented experimental cases where collaborative logistics has been the solution. As part of their study, the authors analyzed the case of the city of Rochelle in France. In this case, they decided to create a distribution center in the outskirts where the providers deliver the merchandise to be classified and organized for further delivery in small electric vehicles. Recently, Makaci et al. (2017) made an empirical study of mutualized distribution centers. The authors studied 7 cases of French companies by carrying out surveys with the people involved in the processes. The authors found that sharing the transportation and storing operations represented a great advantage.

Despite the existence of several studies on truck scheduling problems for CD platforms and mutualization in logistics, to our knowledge, the literature lacks a study on scheduling problems with mutualization in a CD platform. The main contribution of this study is to include the option of leasing the loading and unloading docks of the CD platform for a determined period to allied companies. This option of leasing can be seen as a mutualization operation because allied companies can use a resource of the company and in turn, the company can receive a benefit. Considering the option of leasing the docks change significantly the scheduling optimization problem for CD platforms because it is necessary adding new decision variables and constraints, and modifying the objective function.

3. Truck scheduling problem at a CD platform with mutualized operations

3.1. Problem description

In this section, a formal description of the problem at hand is presented. Let I and O be the set of inbound and outbound trucks from the company owning the platform. Let N and M be the set of loading and unloading docks, respectively, which are independent. Let P be the set of a type of products that circulate through the CD platform and R the set of dock leasing requirements from the allied companies. Each unloading truck $i \in I$ has the following parameters: the arrival time to the platform re_i , the waiting cost at the platform f_i , and the quantity of units that are unloaded at a platform q_{ik} (measured in units) of the product $k \in P$. Similarly, each loading truck $j \in O$ has the following parameters: the arrival time to the platform rs_j , waiting cost at the platform per time unit g_j , and quantity of units that are loaded onto the truck d_{jk} (measured in units) of the product $k \in P$. Considering that the loading and unloading docks are connected by a conveyor belt, the time it takes to transport a type of product unloaded from a truck $i \in I$ until a truck $j \in O$ is constant, and it is defined by the parameter t . Furthermore, the parameter s represents the time to change a truck at the dock. The parameters e and o represent the unloading and loading time per unit, respectively. Each possible requirement $l \in R$ demands an inbound and outbound dock, for a defined period of time within the time interval $[a_l$ and $b_l]$. The income for accepting requirement $l \in R$ is defined as p_l .

In this problem, the objective is to find the order in which the inbound and outbound trucks are served, their assignation to a loading or unloading dock, and to accept or not accept the requirements from allied companies, in a way that the total cost is minimized. The latter is defined as the waiting cost minus the income due to accepting requirements.

To give the reader a better perspective of our problem, we present an illustrative example. For this example, we assume the following sets: $N = \{N1, N2\}$, $M = \{M1, M2, M3\}$, $I = \{IT1, IT2, IT3\}$, $O = \{OT1, OT2, OT3, OT4\}$ and $R = \{R1, R2\}$. Figure 1 shows a solution to the example. In this case, only the requirement $R2$ was accepted, and it was scheduled at the inbound dock $N2$ and outbound dock $M2$. Inbound trucks $IT3$ and $IT1$ were unloaded at unloading dock $N1$. Inbound truck $IT2$ was unloaded at unloading dock $N2$. Outbound trucks $OT4$ and $OT1$ were loaded at dock $M1$. Both trucks needed products from inbound trucks $IT3$ and $IT1$, for that reason their loading operation starts after the unloading operation finishes of trucks $IT3$ and $IT1$. Outbound trucks $OT2$ and $OT3$ were loaded at docks $M3$ and $M2$ respectively. Both trucks needed products from the inbound trucks $IT2$, for that reason their loading operation starts after the unloading operation finishes of the trucks $IT2$.

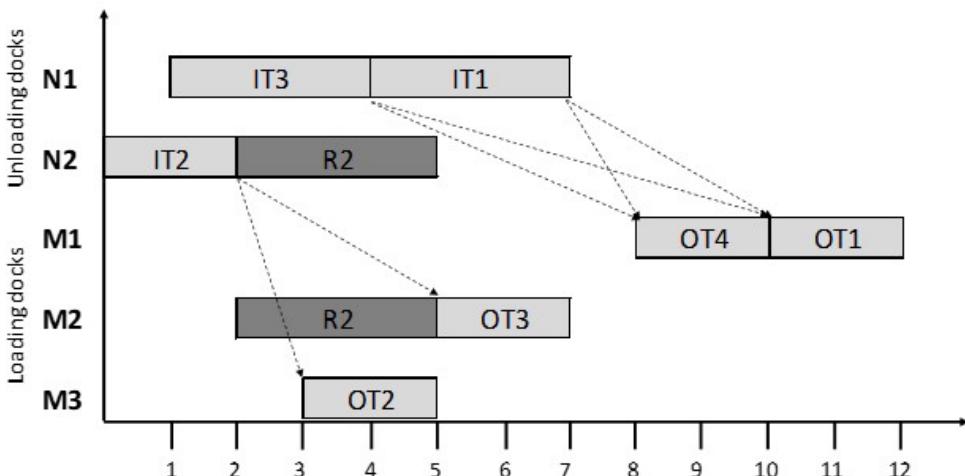


Figure 1. Illustrative example.

3.2. Mathematical model

For the MILP formulation, we introduced the sets $I' = \{0\} \cup I$ y $O' = \{0\} \cup O$, where the element $\{0\}$ represents a dummy truck. Both sets allow defining the precedence decision variables considering that truck $\{0\}$ will always be the first truck served. Additionally, we introduced the parameter $BigM$, which represents a big number. In this mathematical model, we used the following decision variables: variable σ_{ijh} is equal to 1 if the inbound truck $i \in I'$ precedes inbound truck $j \in I'$ at unloading dock $h \in N$ given that $i \neq j$ and 0 otherwise. Variable π_{ijh} is equal to 1 if an outbound truck $i \in O'$ precedes outbound truck $j \in O'$ at loading dock $h \in M$ given that $i \neq j$, and 0 otherwise. Variable λ_{ij} is equal to 1 if any type of product of the inbound truck $i \in I$ is moved to the outbound truck $j \in O$, and 0 otherwise. Variables u_i and v_j represent the end time of the unloading from inbound truck $i \in I'$ and the loading of the outbound truck $j \in O'$. Variables w_i and y_j represent the starting time of the unloading from inbound truck $i \in I'$ and the loading of the outbound truck $j \in O'$ respectively. Variables α_i and β_j represent the waiting time of inbound truck $i \in I$ and the outbound truck $j \in O$ respectively. Variable x_{ijk} is the amount of units moved from inbound truck $i \in I$ to outbound truck $j \in O$ of the product type $k \in P$.

For modeling the acceptance of requirements, the decision variable μ_l is equal to 1 if the requirement $l \in R$ is accepted and 0 otherwise. Moreover, if requirements are accepted, it is necessary to know what loading and unloading dock will be assigned. Variables ε_{lh} and θ_{lk} are equal to 1 if the inbound $h \in N$ and outbound dock $k \in M$ docks are assigned to requirement $l \in R$. On the other hand, it is necessary to be able to control that the accepted requirements do not overlap with any inbound or outbound truck at a given dock (i.e., a requirement and a truck coincide at the same time at the same dock). An alternative for avoiding the overlap is to introduce

a set of time that represents the horizon of the platform schedule. Using that set, we could introduce variables indexed in the time that represent at what moment each activity is carried out. Therefore, the overlapping can be controlled ensuring that almost one activity is performed at a specific moment at a given dock through of a constraint. The main inconvenience of this alternative is the potential impact on the computing time for solving the MILP formulation. This is due to the size of this set of time, which should be significantly large to be able to correctly model the operation. For instance, if the time horizon of the platform is 8 hours and all the associated parameters to time take integer values of minutes, this set of time should have all the minutes in the time horizon of 8 hours (i.e., $\{1, \dots, 480\}$).

As an alternative to avoid the overlap, we use in our model the following decision variables: variable ρ_{ih} is equal to 1 if the inbound truck $i \in I$ is assigned to the unloading dock $h \in N$, and 0 otherwise. Variable ω_{ih} is equal to 1 if the outbound truck $i \in O$ is assigned to the unloading dock $h \in N$, and 0 otherwise. Although those variables can be seen as a subset of the variables σ_{ijh} and π_{ijh} , the motivation to include them is to facilitate the building of the overlapping constraints. In the case that the truck and a requirement are assigned to the same dock, the variable δ_{li} is equal to 1 if the inbound truck $i \in I$ is served first at the dock and then the requirement $l \in R$ without overlapping, otherwise, δ_{li} is equal to 0 if the requirement $l \in R$ is served first at the dock and then the inbound truck $i \in I$ without overlapping. It is important to notice that this variable makes sense only when the requirement $l \in R$ and the inbound truck $i \in I$ are scheduled at the same dock. It means when ε_{lh} and ρ_{ih} take the value of 1. If ε_{lh} or ρ_{ih} are 0, the value of δ_{li} can be neglected. Similarly, the variable δ_{li} functions the binary variable τ_{lj} , but in this case, it is defined for the outbound truck $j \in O$. Considering that overlapping between requirements is a possibility, we introduced the binary variable ψ_{lmh} , which only makes sense if the requirements l and $m \in R$ are scheduled at the same unloading dock $h \in N$. This variable takes a value of 1 if the requirement l is served first and then the requirement m , and it will be 0 if the requirement m is served first and then the requirement l . Finally, the binary variable η_{lmh} is defined for the loading dock $h \in M$ in a similar way to the previous variable. Considering the quantity of sets, parameters and variables, Table 1 summarizes the notation of this problem.

The mathematical model is as follows:

$$\text{Min } Z = \sum_{i \in I} f_i \alpha_i + \sum_{j \in O} g_j \beta_j - \sum_{l \in R} \mu_l p_l \quad (1)$$

The objective function (1) aims to minimize the waiting cost of the inbound and outbound trucks minus the benefits that can be perceived by accepting requirements.

Subject to:

$$\sum_{j \in O} x_{ijk} = q_{ik} \quad \forall i \in I, \forall k \in P \quad (2)$$

$$\sum_{i \in I} x_{ijk} = d_{jk} \quad \forall j \in O, \forall k \in P \quad (3)$$

$$x_{ijk} \leq \text{BigM} \lambda_{ij} \quad \forall i \in I, \forall j \in O, \forall k \in P \quad (4)$$

Constraints (2) control that the inbound trucks are completely unloaded. Constraints (3) control that the outbound trucks are loaded fulfilling with their requirements. Constraints (4) assure for a positive amount of type of product $k \in P$ to be transported between trucks $i \in I$ and $j \in O$, only when $\lambda_{ij} = 1$.

$$u_i \geq w_i + e \sum_{k \in P} q_{ik} \quad \forall i \in I \quad (5)$$

$$v_j \geq y_j + o \sum_{k \in P} d_{jk} \quad \forall j \in O \quad (6)$$

$$w_j \geq u_0 - \text{BigM} (1 - \sigma_{0jh}) \quad \forall j \in I, \forall h \in N \quad (7)$$

Table 1. Variable and parameter definitions.

N	Set of unloading docks
M	Set of loading docks
I	Set of inbound trucks
O	Set of outbound trucks
R	Set of requirements
P	Set of type of products
s	Time to change a truck at the dock
t	Time to transport a type of product unloaded from a truck
$BigM$	Big number
p_l	Income for accepting requirement $l \in R$
a_l	The starting time needed by an inbound and outbound dock for the requirement $l \in R$
b_l	The ending time needed by an inbound and outbound dock for the requirement $l \in R$
re_i, rs_j	Arrival time to the platform of unloading truck $i \in I$ and loading truck $j \in O$
f_{ij}, g_j	Waiting cost at the platform per time unit of unloading truck $i \in I$ and loading truck $j \in O$
q_{ik}	Quantity of units that are unloaded at a platform of the product type $k \in P$ from truck $i \in I$
d_{jk}	Quantity of units that are loaded of the product type $k \in P$ from truck $j \in O$
e	The unloading time per unit
o	The loading time per unit
w_i	The starting time of the unloading from inbound truck $i \in I$
u_i	The end time of the unloading from inbound truck $i \in I$
y_j	The starting time of the loading of outbound truck $j \in O'$
v_j	The end time of the loading of outbound truck $j \in O'$
α_i	The waiting time of inbound truck $i \in I$
β_j	The waiting time of outbound truck $j \in O$
x_{ijk}	The number of units moved from inbound truck $i \in I$ to outbound truck $j \in O$ of the product type $k \in P$
λ_{ij}	1 if any type of product of the inbound truck $i \in I$ is moved to the outbound truck $j \in O$, and 0 otherwise
μ_l	1 if the requirement $l \in R$ is accepted and 0 otherwise
σ_{ijh}	1 if the inbound truck $i \in I'$ precedes inbound truck $j \in I'$ at unloading dock $h \in N$ given that $i \neq j$ and 0 otherwise
π_{ijh}	1 if an outbound truck $i \in O'$ precedes outbound truck $j \in O'$ at loading dock $h \in M$ given that $i \neq j$, and 0 otherwise
$\varepsilon_{lh}, \theta_{lk}$	1 if the inbound $h \in N$ and outbound dock $k \in M$ docks are assigned to requirement $l \in R$
δ_{li}	1 if the inbound truck $i \in I$ is served first at the dock and then the requirement $l \in R$ without overlapping
τ_{lj}	1 if the outbound truck $j \in O$ is served first at the dock and then the requirement $l \in R$ without overlapping
ρ_{ih}	1 if the truck $i \in I$ is assigned to the loading dock $h \in N$
ω_{ih}	1 if the truck $i \in O$ is assigned to the loading dock $h \in M$
ψ_{lmh}	Which only makes sense if the requirements l and $m \in R$ are scheduled at the same unloading dock $h \in N$
η_{lmh}	Which only makes sense if the requirements l and $m \in R$ are scheduled at the same loading dock $h \in M$

$$y_j \geq v_0 - BigM(1 - \pi_{0jh}) \quad \forall j \in O, \forall h \in M \quad (8)$$

$$w_j \geq u_i + s - BigM(1 - \sigma_{ijh}) \quad \forall i, j \in I, \forall h \in N : i \neq j \quad (9)$$

$$y_j \geq v_i + s - BigM(1 - \pi_{ijh}) \quad \forall i, j \in O, \forall h \in M : i \neq j \quad (10)$$

$$w_j \geq re_i \quad \forall j \in I \quad (11)$$

$$y_j \geq rs_j \quad \forall j \in O \quad (12)$$

$$y_j \geq u_i + t - BigM(1 - \lambda_{ij}) \quad \forall i \in I, \forall j \in O \quad (13)$$

$$\alpha_i \geq u_i - re_i \quad \forall i \in I \quad (14)$$

$$\beta_j \geq v_j - rs_j \quad \forall j \in O \quad (15)$$

Constraints (5) and (6) define the end time for the unloading and loading of the inbound and outbound trucks, respectively. Constraints (7) and (8) define the starting time of the unloading and loading of the first truck at each dock. Constraints (9) and (10) define the starting time of the unloading and loading for the rest of the trucks (*i.e.*, from the second truck onwards) at each dock. It is important to notice that the first vehicle that arrives at the dock does not incur in the time s associated with the change of truck. Constraints (11) and (12) guarantee that the starting time for loading and unloading for the trucks is higher than their respective arrival time. Constraints (13) control for the starting time of an outbound truck $j \in O$ to be higher than the ending time of the unloading of the inbound vehicle $i \in I$ plus the merchandise's moving time, given that any merchandise is moved between the trucks $i \in I$ and $j \in O$. Constraints (14) and (15) define the waiting time of the inbound and outbound trucks.

$$\sum_{i \in I', i \neq j} \sum_{h \in N} \sigma_{ijh} = 1 \quad \forall j \in I \quad (16)$$

$$\sum_{i \in O', i \neq j} \sum_{h \in M} \pi_{ijh} = 1 \quad \forall j \in O \quad (17)$$

$$\sum_{j \in I', i \neq j} \sigma_{ijh} - \sum_{j \in I', i \neq j} \sigma_{jih} = 0 \quad \forall i \in I', \forall h \in N \quad (18)$$

$$\sum_{j \in O', i \neq j} \pi_{ijh} - \sum_{j \in O', i \neq j} \pi_{jih} = 0 \quad \forall i \in O', \forall h \in M \quad (19)$$

$$\sum_{j \in I'} \sigma_{0jh} \leq 1 \quad \forall h \in N \quad (20)$$

$$\sum_{j \in O'} \pi_{0jh} \leq 1 \quad \forall h \in M \quad (21)$$

Constraints (16) and (17) assure that all the inbound and outbound trucks are served. Constraints (18) and (19) impose the preservation of flow for the variable σ_{ijh} and π_{ijh} . Constraints (20) and (21) control for every loading and unloading dock to be assigned one vehicle at most. The later constraints with the constraints (18) and (19) control that only one truck is served at a time at a dock.

$$\mu_l = \sum_{h \in N} \varepsilon_{lh} \quad \forall l \in R \quad (22)$$

$$\mu_l = \sum_{h \in M} \theta_{lh} \quad \forall l \in R \quad (23)$$

$$\sum_{h \in N} \varepsilon_{lh} \leq 1 \quad \forall l \in R \quad (24)$$

$$\sum_{h \in M} \theta_{lh} \leq 1 \quad \forall l \in R \quad (25)$$

$$\rho_{ih} \geq \sum_{j \in I'} \sigma_{ijh} \quad \forall i \in I, \forall h \in N \quad (26)$$

$$\sum_{h \in N} \rho_{ih} = 1 \quad \forall i \in I \quad (27)$$

$$\omega_{ih} \geq \sum_{j \in O'} \pi_{ijh} \quad \forall i \in O, \forall h \in M \quad (28)$$

$$\sum_{h \in M} \omega_{ih} = 1 \quad \forall i \in O \quad (29)$$

Constraints (22) and (23) assure that a requirement must be scheduled at an inbound and outbound dock when requirements are accepted. Constraints (24) and (25) control for a requirement to be assigned to one inbound and outbound dock at most. Constraints (26) assure that the inbound truck ρ_{ih} takes a value of 1 if the truck $i \in I$ is assigned to the loading dock $h \in N$. Constraints (27) assure that each inbound truck to be only assigned to one unloading dock. Constraints (28) assure that the inbound truck ω_{ih} takes a value of 1 if the truck $i \in O$ is assigned to the loading dock $h \in M$. Constraints (29) assure that each outbound truck to be only assigned to one loading dock.

For understanding the following constraints, Figure 2 shows the possible cases when an inbound truck $i \in I$ and a requirement $l \in R$ are assigned to the same dock $h \in N$. Subfigures 2a and 2b present the feasible cases (i.e., there is no overlapping). When the inbound truck $i \in I$ is served first than the requirement $l \in R$ (Figure 2a) and when the requirement $l \in R$ is ahead of the inbound truck $i \in I$ (Figure 2b). Subfigures 2c and 2d present the unfeasible cases (i.e., when there is overlapping). When the truck $i \in I$ is served first and before finishing requirement $l \in R$ starts to be served (Figure 2c), and when the requirement $l \in R$ starts first and before finishing the truck $i \in I$ starts to be served (Figure 2d).

$$a_l - u_i \geq \text{BigM}(\delta_{li} - 1) - \text{BigM}(2 - \varepsilon_{lh} - \rho_{ih}) \quad \forall i \in I, \forall l \in R, \forall h \in N \quad (30)$$

$$b_l - w_i \geq \text{BigM}(\delta_{li} - 1) - \text{BigM}(2 - \varepsilon_{lh} - \rho_{ih}) \quad \forall i \in I, \forall l \in R, \forall h \in N \quad (31)$$

$$a_l - u_i \leq \text{BigM}\delta_{li} + \text{BigM}(2 - \varepsilon_{lh} - \rho_{ih}) \quad \forall i \in I, \forall l \in R, \forall h \in N \quad (32)$$

$$b_l - w_i \leq \text{BigM}\delta_{li} + \text{BigM}(2 - \varepsilon_{lh} - \rho_{ih}) \quad \forall i \in I, \forall l \in R, \forall h \in N \quad (33)$$

Constraints (30-33) guarantee that the inbound truck $i \in I$ and a requirement $l \in R$ will not overlap if they are assigned to the same dock $h \in N$. In these constraints, the start of the requirement a_l is always compared against the end of the service of the truck u_i , and the end of the requirement b_l against the start of the service of the truck w_i , as presented in Figure 2. It is important to notice that the constraints will be trivially met if $\varepsilon_{lh} = 0$ (i.e., the requirement $l \in R$ is not assigned to the dock $h \in N$) or $\rho_{ih} = 0$ (i.e., the truck $i \in I$ is not assigned to the dock $h \in N$). However, in the case that $\varepsilon_{lh} = 1$ and $\rho_{ih} = 1$, for obtaining a feasible solution the inbound truck $i \in I$ and a requirement $l \in R$ must be scheduled as the shown cases Figures 2a or 2b. For meeting the case in Figure 2a (i.e., $a_l - u_i \geq 0$, and at the same time $b_l - w_i \geq 0$) δ_{li} must take the value of 1. For meeting the case in Figure 2b (i.e., $a_l - u_i \leq 0$, and at the same time $b_l - w_i \leq 0$) δ_{li} must take the value of 0. Therefore, the feasible case presented in Figure 2a will occur for $\delta_{li} = 1$, and the feasible case presented in Figure 2b will occur for $\delta_{li} = 0$. It is important to notice that with these constraints, the unfeasible cases presented in Figures 2c and 2d will not take place.

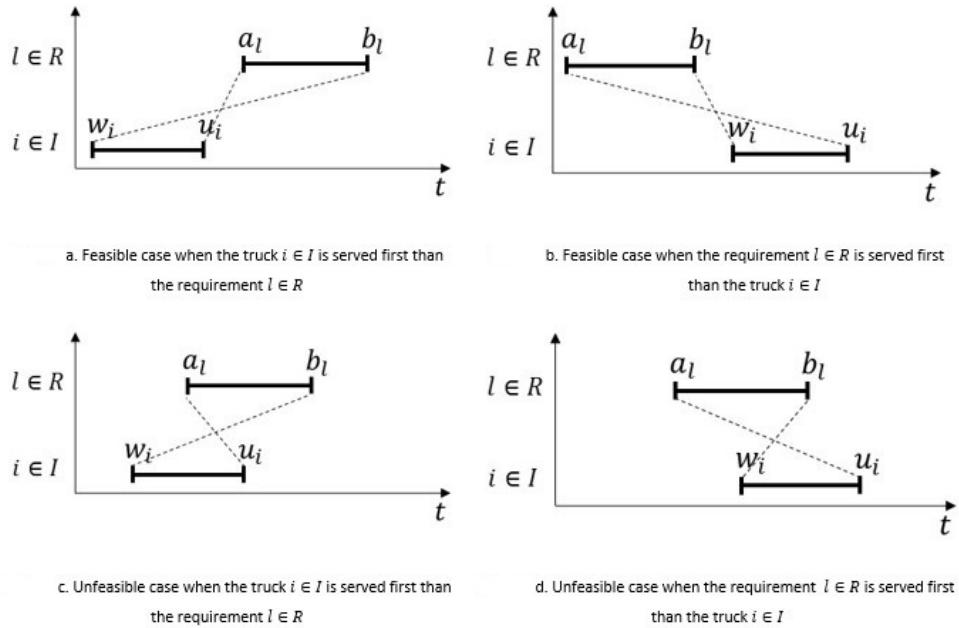


Figure 2. Possible cases of scheduling when an inbound truck $i \in I$ and a requirement $l \in R$ are assigned to the same dock $h \in N$.

$$a_l - v_j \geq \text{BigM}(\tau_{li} - 1) - \text{BigM}(2 - \theta_{lh} - \omega_{ih}) \quad \forall i \in O, \forall l \in R, \forall h \in M \quad (34)$$

$$b_l - y_i \geq \text{BigM}(\tau_{li} - 1) - \text{BigM}(2 - \theta_{lh} - \omega_{ih}) \quad \forall i \in O, \forall l \in R, \forall h \in M \quad (35)$$

$$a_l - v_j \leq \text{BigM}\tau_{li} + \text{BigM}(2 - \theta_{lh} - \omega_{ih}) \quad \forall i \in O, \forall l \in R, \forall h \in M \quad (36)$$

$$b_l - y_i \leq \text{BigM}\tau_{li} + \text{BigM}(2 - \theta_{lh} - \omega_{ih}) \quad \forall i \in O, \forall l \in R, \forall h \in M \quad (37)$$

Constraints (34-37) guarantee that the outbound truck $i \in O$ and a requirement $l \in R$ will not overlap if they are assigned to the same dock $h \in M$. These constraints work similarly to constraints (30-33).

$$a_l - b_m \geq \text{BigM}(\psi_{lmh} - 1) - \text{BigM}(2 - \varepsilon_{lh} - \varepsilon_{mh}) \quad \forall l, m \in R, \forall h \in N, l \neq m \quad (38)$$

$$b_l - a_m \geq \text{BigM}(\psi_{lmh} - 1) - \text{BigM}(2 - \varepsilon_{lh} - \varepsilon_{mh}) \quad \forall l, m \in R, \forall h \in N, l \neq m \quad (39)$$

$$a_l - b_m \leq \text{BigM}\psi_{lmh} + \text{BigM}(2 - \varepsilon_{lh} - \varepsilon_{mh}) \quad \forall l, m \in R, \forall h \in N, l \neq m \quad (40)$$

$$b_l - a_m \leq \text{BigM}\psi_{lmh} + \text{BigM}(2 - \varepsilon_{lh} - \varepsilon_{mh}) \quad \forall l, m \in R, \forall h \in N, l \neq m \quad (41)$$

$$a_l - b_m \geq \text{BigM}(\eta_{lmk} - 1) - \text{BigM}(2 - \theta_{lk} - \theta_{mk}) \quad \forall l, m \in R, \forall k \in M, l \neq m \quad (42)$$

$$b_l - a_m \geq \text{BigM}(\eta_{lmk} - 1) - \text{BigM}(2 - \theta_{lk} - \theta_{mk}) \quad \forall l, m \in R, \forall k \in M, l \neq m \quad (43)$$

$$a_l - b_m \leq \text{BigM} \eta_{lmk} + \text{BigM} (2 - \theta_{lk} - \theta_{mk}) \quad \forall l, m \in R, \forall k \in M, l \neq m \quad (44)$$

$$b_l - a_m \leq \text{BigM} \eta_{lmk} + \text{BigM} (2 - \theta_{lk} - \theta_{mk}) \quad \forall l, m \in R, \forall k \in M, l \neq m \quad (45)$$

Considering that there are possibilities for the requirements l and $m \in R$ to overlap at the inbound dock $h \in N$ and outbound dock $k \in M$, constraints (38-41) guarantee that the requirements will not overlap with one another at an inbound dock. Similarly, constraints (42-45) guarantee that the requirements will not overlap with one another at an outbound dock. These constraints work similarly to constraints (30-33).

$$w_i, u_i, \alpha_i \geq 0 \quad \forall i \in I \quad (46)$$

$$y_j, v_j, \beta_j \geq 0 \quad \forall j \in O \quad (47)$$

$$x_{ijk} \in Z^+ \quad \forall i \in I, \forall j \in O, \forall k \in P \quad (48)$$

$$\lambda_{ij} \in \{0,1\} \quad \forall i \in I, \forall j \in O \quad (49)$$

$$\rho_{ih} \in \{0,1\} \quad \forall i \in I, \forall h \in N \quad (50)$$

$$\omega_{ih} \in \{0,1\} \quad \forall i \in O, \forall h \in M \quad (51)$$

$$\sigma_{ijh} \in \{0,1\} \quad \forall i, j \in I', \forall h \in N \quad (52)$$

$$\pi_{ijh} \in \{0,1\} \quad \forall i, j \in O', \forall h \in N \quad (53)$$

$$\varepsilon_{lh} \in \{0,1\} \quad \forall l \in R, \forall h \in N \quad (54)$$

$$\theta_{lh} \in \{0,1\} \quad \forall l \in R, \forall h \in M \quad (55)$$

$$\delta_{li} \in \{0,1\} \quad \forall l \in R, \forall i \in N \quad (56)$$

$$\tau_{li} \in \{0,1\} \quad \forall l \in R, \forall i \in O \quad (57)$$

$$\psi_{lmh} \in \{0,1\} \quad \forall l, m \in R, \forall h \in N \quad (58)$$

$$\eta_{lmh} \in \{0,1\} \quad \forall l, m \in R, \forall h \in M \quad (59)$$

Constraints (46 – 59) define the domain of the decision variables.

4. Computational experiments

Aiming to evaluate the performance of the MILP formulation in solving the studied problem, we used a set of instances based on a real operation from a retail company. This company operates under a CD operation, and it has the capability of leasing some of the docks in its platform to allied companies. In this experiment, we analyzed the behavior of the MILP formulation performance in terms of the different characteristics of the problem. The MILP formulation was solved using Gurobi 8.1.1 with a time limit of 10,800s for each instance. All the experiments were run on a computing cluster with Intel Xeon E5-2683 v4 processor (32 cores at 2.1 GHz) and 64 GB of RAM, both running on Linux Rocks 6.2-64 bits – Centos 6.6.

4.1. Test instances

We generated a set of 128 instances inspired by the real operation of a retail company. Aiming to obtain different size instances, we generated instances with different amounts of unloading $|N|$ and loading $|M|$ docks, inbound $|I|$ and outbound $|O|$ trucks, requirements $|R|$, and product types $|P|$. Table 2 contains the values of the parameters used in the generation of the instances. The magnitude of the used parameters is based on the information from the real operation. Using these values, we obtained 64 combinations. From each combination, we generated 2 instances; hence, obtaining 128 instances. The magnitudes of arrival time of the trucks, vehicle change time, moving time between loading and unloading docks, loading and unloading rate, and the amount of units per product were randomly generated using magnitudes from the real operation.

Table 2. Parameters for the generation of the instances.

Parameter	Nomenclature	Value
Unloading docks	$ N $	{2, 4}
Loading docks	$ M $	{4, 6}
Inbound trucks	$ I $	{4, 8}
Outbound trucks	$ O $	{8, 16}
Requirements	$ R $	{10, 20}
Product types	$ P $	{20, 40}

4.2. Results

In this section, we analyze the computational results about the performance of the MILP formulation. The detailed results for each instance are documented in Table A1. Table 3 shows the results consolidated from the instances wherein the model reached the optimal solution. Row 1 presents how many instances were optimally solved, and rows 2-5 present the information concerning average, maximum, minimum computational times, and the standard deviation (in seconds), respectively. Results in Table 3 demonstrate that the MILP formulation is capable of finding the optimal solution for 65/128 instances. The maximum CPU time of these 65 instances was 8103 seconds, which is a reasonable time and meets the expectations of the operation for the company that inspired this study.

Table 3. Optimal solution results.

Metric	Value
Number of optimal solutions	65
Average CPU time (sec.)	329.75
Maximum CPU time (sec.)	8103
Minimum CPU time (sec.)	0.86
Standard deviation (sec)	1,069

Figure 3 presents the behavior of the number of optimal solutions found within the time. As seen in Figure 3, we observe that 91% of the optimal solutions have been found within 958 seconds, which also allows us to conclude that the MILP formulation behaves in an acceptable way.

Table 4 shows the results for the instances that did not find an optimal solution (63/128). As seen in Table 4, the average optimality gap is 8.95%, the maximum gap is 38%, the minimum gap is 0.06%, and the standard deviation is 7.9%. Analyzing the results in detail, we found that 25 of the instances wherein an optimal solution was not reached have a gap lesser than 5%. Despite not finding the optimal solution in these instances, the behavior of the gap is acceptable for the problem at hand.

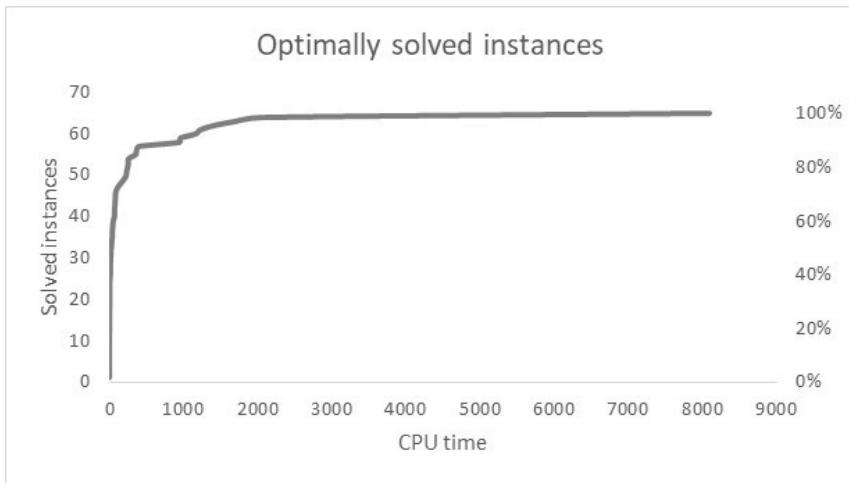


Figure 3. Number of optimal solutions found within time.

Table 4. Results of the instances that did not reach an optimal solution.

Metric	Value
Average optimality gap	8.95%
Maximum optimality gap	38.38%
Minimum optimality gap	0.06%
Standard deviation (%)	7.9%

Hereafter, we analyze the behavior of the optimality gap according to the characteristics of the problem. Figures 4 and 5 present the optimality gap for $|P|=20$ and $|P|=40$, respectively. In both figures, the X axis presents the loading docks and the outbound trucks; the Y axis presents the loading docks and the inbound trucks. The gray and black colored circles represent the problems with $|R|=10$ and $|R|=20$, respectively. The size of the circles represents the value of the optimality gap.

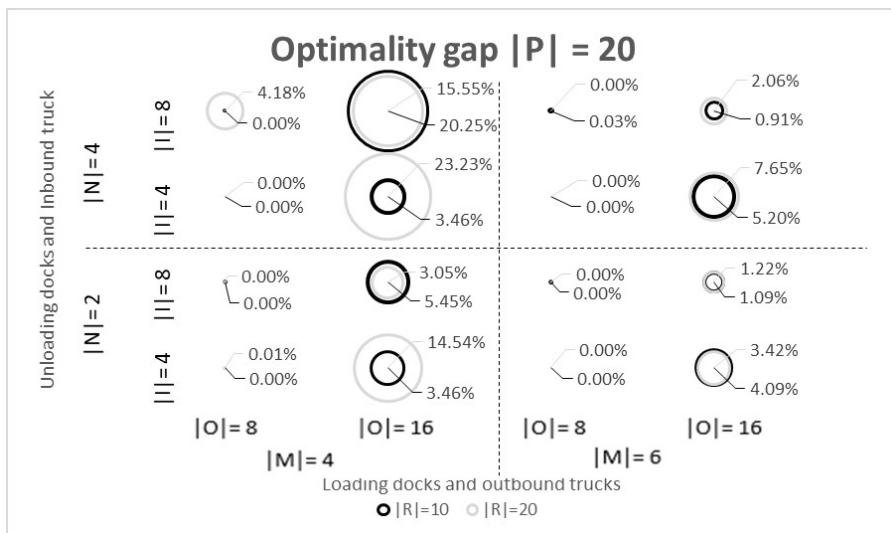


Figure 4. Behavior of the optimality gap as a function of the characteristics of the instances with $|P|=20$.

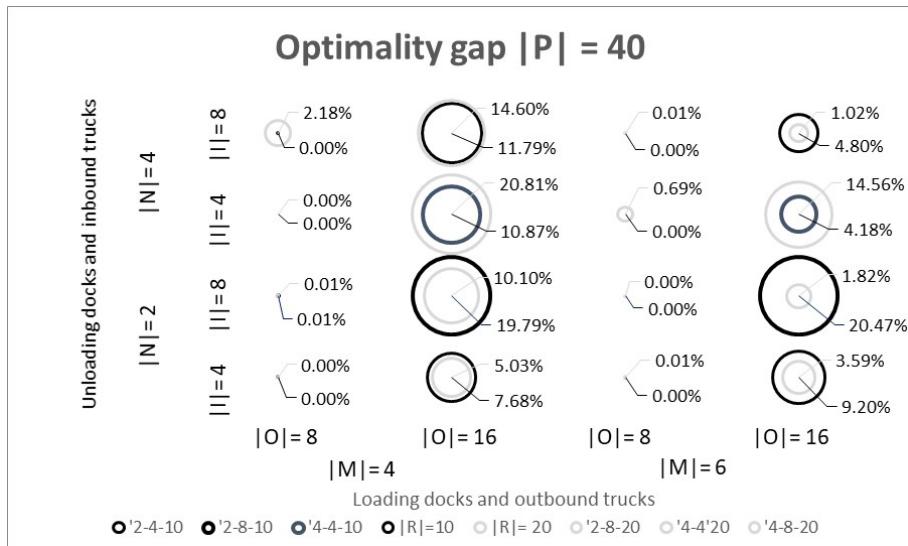


Figure 5. Behavior of the optimality gap as a function of the characteristics of the instances with $|P|=40$.

After comparing Figures 4 and 5, there is no notable impact of the optimality gap behavior in terms of the number of product types. Nonetheless, the average gap for the instances $|P|=20$ and $|P|=40$ is 3.71% and 5.10%, respectively. This result shows that, in average, the instances with 40 products are more difficult to solve than the instances with 20 products. As expected, there is a low optimality gap when the number of requirements is $|R|=10$ when compared with $|R|=20$. Conversely, the results do not show that there is a low optimality gap value when there are less inbound or outbound docks. Both Figures 4 and 5 show cases wherein a higher number of outbound docks have a lower optimality gap value. For instance, in Figure 4, when the instances have 6 loading docks, the optimality gap is generally lower than when the instances have 4 loading docks. Figure 5 shows that for instances $|I|=8$ and $|O|=16$, the optimality gap is lower when $|N|=4$ than $|N|=2$. This behavior, having a higher optimality gap for fewer docks, can be attributed to the fact that the problem becomes more constrained; hence, more difficult to solve. Regarding the number of outbound trucks, the results show that the optimality gap tends to increase when there are more outbound trucks. Nonetheless, this phenomenon is not evident with the inbound trucks. In some cases, the more inbound trucks there are, the better the optimality gap. This case can be seen in Figure 5 with $|M|=6$, $|N|=4$, and $|O|=16$, where the optimality gap is lower for $|I|=8$ than $|I|=4$.

Lastly, we analyzed the solutions in terms of the accepted requirements. The solutions showed that all the solutions accept requirements. We found that, in average, 48% of the requirements are accepted, and in one instance, all the requirements were accepted. These results show that, for the operation that inspired this problem, it is interesting to lease the docks to allied companies as a way to subsidize the operational costs.

5. Conclusions

This study proposes a new scheduling problem for a CD platform. This problem aims to find the scheduling of both inbound and outbound trucks for the company owning the platform while deciding whether or not to accept requirements to lease their docks to allied companies. In the case of accepting the requirements, it has to be decided which docks will be assigned. These decisions have to be made in a way that minimize the operational cost minus an income due to the leasing.

For solving this problem, we propose a MILP formulation. To test this MILP formulation, we built a set of 128 instances, which are based on a real CD operation. Through a computational experiment, we found that the MILP formulation has an acceptable behavior. Using a commercial solver, the MILP formulation is capable of finding 65/128 optimal solutions with an average time of 4.41%. For the remaining 63/128 instances, the MILP formulation obtains an optimality gap of 8.95% in 10800 seconds. For the real case that inspired this study, these results are acceptable given that the scheduling of the trucks is done during the night hours of the day prior to the operation. When analyzing the solutions, we found that 128/128 instances accepted the requirements. Therefore, we can conclude that, for the company that inspired this problem, it is economically interesting to lease some of its docks to allied companies.

Interesting research directions include considering other resources of the CD platform for leasing such as people and lift trucks. Another interesting research direction is to include in this problem the routing decisions of the trucks, especially those of the outbound trucks.

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Appendix 1. Detailed results.

Table A1 presents the detailed results for the computational experiment from the MILP formulation. Column 1 presents a consecutive number for each of the instances, columns 2 – 7 present the parameters for the generated instances, column 8 shows the optimality gap, column 9 presents the CPU time, and column 10 presents the objective function value.

Table A1. Detailed results.

Instance ID	Parameters						Results		
	I :	O :	P :	N :	M :	R :	Gap	CPU time (s)	Objective function
1	4	8	20	2	4	10	0.00%	1.24	312,858.00
2	4	8	20	2	4	20	0.01%	70.95	255,583.00
3	4	8	20	2	6	10	0.00%	5.92	564,898.00
4	4	8	20	2	6	20	0.00%	15.42	345,280.00
5	4	8	20	4	4	10	0.00%	1.15	865,803.00
6	4	8	20	4	4	20	0.00%	5.34	138,327.00
7	4	8	20	4	6	10	0.00%	0.91	492,354.00
8	4	8	20	4	6	20	0.00%	6.52	328,345.00
9	4	8	40	2	4	10	0.00%	2.78	708,240.00
10	4	8	40	2	4	20	0.00%	5.4	629,168.00
11	4	8	40	2	6	10	0.00%	5.59	591,834.00
12	4	8	40	2	6	20	0.01%	6.35	603,579.00
13	4	8	40	4	4	10	0.00%	1.66	573,601.00
14	4	8	40	4	4	10	0.00%	0.86	406,284.00
15	4	8	40	4	6	10	0.00%	1.2	935,785.00
16	4	8	40	4	6	20	1.38%	10800	652,717.00
17	4	16	20	2	4	10	1.05%	10800	835,855.00
18	4	16	20	2	4	20	12.40%	10800	835,566.00
19	4	16	20	2	6	10	8.18%	10800	771,120.00
20	4	16	20	2	6	20	0.00%	249.25	380,905.00
21	4	16	20	4	4	10	6.91%	10800	719,676.00
22	4	16	20	4	4	20	19.18%	10800	951,634.00
23	4	16	20	4	6	10	10.38%	10800	187,914.00
24	4	16	20	4	4	20	27.28%	10800	690,746.00
25	4	16	40	2	4	10	9.04%	10800	1,428,986.00
26	4	16	40	2	4	20	4.24%	10800	1,783,752.00
27	4	16	40	2	6	10	7.52%	10800	912,581.00
28	4	16	40	2	6	20	4.86%	10800	1,095,480.00
29	4	16	40	4	4	10	10.81%	10800	819,178.00
30	4	16	40	4	4	20	3.23%	10800	1,194,469.00
31	4	16	40	4	6	10	5.61%	10800	1,437,353.00
32	4	16	40	4	6	10	2.75%	10800	1,202,794.00
33	8	8	20	2	4	10	0.01%	70.98	844,308.00
34	8	8	20	2	4	20	0.00%	23.96	546,287.00
35	8	8	20	2	6	10	0.00%	88.46	917,993.00
36	8	8	20	2	6	20	0.00%	41.25	673,694.00
37	8	8	20	4	4	10	0.01%	2054.39	497,865.00
38	8	8	20	4	4	20	0.00%	1404.75	350,849.00
39	8	8	20	4	6	10	0.06%	10800	294,047.00
40	8	8	20	4	6	20	0.00%	80.42	642,000.00
41	8	8	40	2	4	10	0.00%	368.2	1,058,835.00
42	8	8	40	2	4	10	0.01%	1161.51	886,919.00
43	8	8	40	2	6	10	0.00%	15.93	1,135,549.00
44	8	8	40	2	6	20	0.00%	22.66	1,142,830.00
45	8	8	40	4	4	10	0.00%	33.28	544,765.00
46	8	8	40	4	4	20	4.36%	10800	927,428.00
47	8	8	40	4	6	10	0.00%	39.96	464,805.00
48	8	8	40	4	6	20	0.00%	154.18	1,030,667.00

Table A1. Continued...

Instance ID	Parameters						Results		
	I :	O :	P :	N :	M :	R :	Gap	CPU time (s)	Objective function
49	8	16	20	2	4	10	5.91%	10800	1,488,440.00
50	8	16	20	2	4	20	3.64%	10800	911,947.00
51	8	16	20	2	6	10	0.60%	10800	1,334,189.00
52	8	16	20	2	6	20	2.44%	10800	1,040,312.00
53	8	16	20	4	4	10	19.45%	10800	810,105.00
54	8	8	20	4	6	20	0.00%	110.78	1,073,301.00
55	8	16	20	4	6	10	1.62%	10800	533,236.00
56	8	16	20	4	4	20	19.72%	10800	542,777.00
57	8	16	40	2	4	10	27.40%	10800	2,292,650.00
58	8	16	40	2	4	20	9.87%	10800	1,787,240.00
59	8	16	40	2	6	10	13.60%	10800	2,066,403.00
60	8	16	40	2	6	20	0.41%	10800	1,414,612.00
61	8	16	40	4	4	10	16.10%	10800	1,763,647.00
62	8	16	40	4	4	20	11.98%	10800	1,337,360.00
63	8	16	40	4	6	10	4.81%	10800	1,218,799.00
64	8	16	40	4	6	20	0.62%	10800	868,508.00
65	4	8	20	2	4	10	0.00%	1.06	505,126.00
66	4	8	20	2	4	20	0.00%	8.65	646,574.00
67	4	8	20	2	6	10	0.00%	2.22	527,179.00
68	4	8	20	2	6	20	0.00%	16.56	524,141.00
69	4	8	20	4	4	10	0.00%	5.06	202,053.00
70	4	8	20	4	4	20	0.00%	32.48	301,253.00
71	4	8	20	4	6	10	0.00%	2.28	219,005.00
72	4	8	20	4	6	20	0.00%	19.17	428,856.00
73	4	8	40	2	4	10	0.00%	77.35	127,437.00
74	4	8	40	2	4	20	0.01%	198.05	815,770.00
75	4	8	40	2	6	10	0.00%	3.24	701,028.00
76	4	8	40	2	6	20	0.00%	5.86	583,738.00
77	8	8	40	4	6	20	0.01%	1692.53	426,410.00
78	4	8	40	4	4	20	0.00%	4.3	500,186.00
79	4	8	40	4	6	10	0.00%	1.97	660,610.00
80	4	8	40	4	6	20	0.00%	3.28	211,708.00
81	4	16	20	2	4	10	5.87%	10800	1,796,690.00
82	4	16	20	2	4	20	16.68%	10800	585,802.00
83	4	16	20	2	6	10	0.01%	234.26	832,645.00
84	4	16	20	2	6	20	6.85%	10800	623,745.00
85	4	16	20	4	4	10	0.01%	8103.08	844,798.00
86	8	16	20	4	6	20	3.16%	10800	742,344.00
87	4	16	20	4	6	10	0.01%	958.22	541,751.00
88	8	16	20	4	6	20	0.95%	10800	788,795.00
89	4	16	40	2	4	10	6.33%	10800	490,816.00
90	4	16	40	2	4	20	5.83%	10800	892,754.00
91	4	16	40	2	6	10	10.89%	10800	1,044,698.00
92	4	16	40	2	6	20	2.31%	10800	751,663.00
93	4	16	40	4	4	10	10.93%	10800	1,550,992.00
94	4	16	40	4	4	20	38.38%	10800	344,941.00
95	4	16	20	4	6	20	7.53%	10800	736,943.00
96	4	16	40	4	6	20	11.85%	10800	551,485.00
97	8	8	20	2	4	10	0.00%	405.39	513,908.00
98	8	8	20	2	4	20	0.01%	258.47	606,950.00
99	8	8	20	2	6	10	0.00%	228.94	1,014,807.00
100	8	8	20	2	6	20	0.00%	48.66	608,272.00
101	8	8	20	4	4	10	0.00%	5.85	299,956.00
102	8	8	20	4	4	20	8.37%	10800	405,949.00

Table A1. Continued...

Instance ID	Parameters						Results		
	I :	O :	P :	N :	M :	R :	Gap	CPU time (s)	Objective function
103	8	8	20	4	6	10	0.00%	14.93	466,094.00
104	4	16	20	4	6	20	7.78%	10800	1,293,690.00
105	4	16	40	4	6	20	17.27%	10800	649,728.00
106	8	8	40	2	4	20	0.01%	53.75	461,588.00
107	8	8	40	2	6	10	0.00%	43.23	855,876.00
108	8	8	40	2	6	20	0.01%	1232.06	813,503.00
109	8	8	40	4	4	10	0.00%	72.47	547,090.00
110	8	8	40	4	4	20	0.01%	945.76	457,167.00
111	8	8	40	4	6	10	0.00%	261.11	879,908.00
112	8	8	40	2	4	20	0.01%	354.82	1,101,907.00
113	8	16	20	2	4	10	4.99%	10800	1,129,864.00
114	8	16	20	2	4	20	2.47%	10800	1,525,942.00
115	8	16	20	2	6	10	1.59%	10800	906,274.00
116	8	16	20	2	6	20	0.00%	83.88	652,993.00
117	8	16	20	4	4	10	21.05%	10800	892,154.00
118	8	16	20	4	4	20	11.37%	10800	745,172.00
119	8	16	20	4	6	10	0.20%	10800	823,104.00
120	8	16	40	4	6	20	1.42%	10800	706,096.00
121	8	16	40	2	4	10	12.19%	10800	1,754,913.00
122	8	16	40	2	4	20	10.33%	10800	1,239,179.00
123	8	16	40	2	6	10	27.34%	10800	2,007,234.00
124	8	16	40	2	6	20	3.23%	10800	1,257,199.00
125	8	16	40	4	4	10	7.48%	10800	1,133,588.00
126	8	16	40	4	4	20	17.23%	10800	1,215,819.00
127	8	16	40	4	6	10	4.79%	10800	1,164,960.00
128	4	8	40	4	4	20	0.00%	3.68	716,780.00